

Comp Sci 212

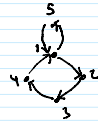
1. Trees

Theorem. The following are equivalent (and define trees)

1. $G = (V, E)$ is connected and acyclic (no subgraph)
2. Every pair of vertices is connected by a unique path (no repeats)
3. G is connected and $|E| = |V| - 1$.
4. G is acyclic and $|E| = |V| - 1$
5. G is acyclic and adding any new edge creates a cycle.



Proof. Show that $1 \rightarrow 2, 2 \rightarrow 3, 3 \rightarrow 4, 4 \rightarrow 1, 5 \rightarrow 1, 1 \rightarrow 5$

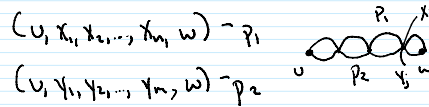


$4 \rightarrow 1$ (Friday)

$1 \rightarrow 2$ Use contradiction

Let u, w be two vertices $\rightarrow$ not connected, then G is not connected

Assume there's two paths connecting $u \leftrightarrow w$



Let i be the largest number s.t. x_i is in P_2 , as y_j . (if no such x_i , use u)

Then $(w, x_n, x_{n-1}, \dots, x_i, y_j, y_{j-1}, \dots, y_1, u)$ is a cycle.

$2 \rightarrow 3$ Strong induction on $|V|$

Base case - $|V| = 1$, obvious

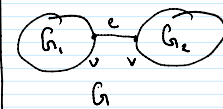
Inductive step - Assume true if $|V| \leq k$

If $|V| = k+1$, let $\{u, v\} \in E$ let $G' = (V, E \setminus \{e\})$

In G' , $u \leftrightarrow v$ must be disconnected, otherwise there would be 2 paths from u to v in G .

Then G' has two connected components, $G_1 + G_2$

$G_i = (V_i, E_i)$



$G_1 + G_2$ satisfy 2, so $|E_i| \leq |V_i| - 1$ for $i=1, 2$

$|E| = |E_1| + |E_2| + 1 \leq |V_1| - 1 + |V_2| - 1 + 1 = |V| - 1$

$|V| = |V_1| + |V_2|$

$3 \rightarrow 4$ Use contradiction

(sketch) Assume G contains a cycle, k edges, k vertices.

Add vertices one by one. Adding each vertex must involve adding an edge, because graph is connected. Ends with $|E| \geq |V|$, contradicts the assumption that $|E| = |V| - 1$.

Each time you add a vertex, choose a vertex that neighbors a vertex already in the graph. Such a vertex must exist, otherwise the graph is not connected. Let G' be graph that's built so far, let v be a vertex in G' but not G , w be a vertex in G ($v, x_1, x_2, \dots, x_n, w$). Let i be the largest s.t.

$5 \rightarrow 1$ Use contradiction. Assume G is not connected.

In particular, $u \leftrightarrow v$ be two vertices not connected.

Adding $\{u, v\}$ to E can not create a cycle. Contradicts 5.

$1 \rightarrow 5$ If we add $\{u, v\}$ there is a path from u to v , $(u, v_1, v_2, \dots, v_n, v)$, adding $\{u, v\}$ turns path into a cycle.

x_i is not in G' . Then adding x_i to G' must add an edge, $\{x_i, x_{i+1}\}$ ($\{x_i, w\}$ if $i=n$).