

Boolean Functions

1. Summary
 - Switching Lemma
 - LMN Theorem
2. Q+A
3. Problems

Announcements

- Theory slack / seminar (Friday)
- TCS+ February 17 (Circuit Lower Bounds from Pseudorandom Restrictions)

LMN - size s depth- d circuit, spectrum is ϵ -concentrated up to degree $O(\log(\frac{s}{\epsilon})^{d-1} \cdot \log(\frac{1}{\epsilon}))$.

Corollary - $\text{poly}(n)$ -size circuit, depth- d , can be learned up to $\text{poly}(\epsilon)$ error in time $n^{O(\log n)^d}$.

Corollary - If C is a depth- d circuit st $\Pr[C(x) = \chi_{\text{maj}}(x)] \geq \frac{1}{2} + \epsilon_0$, size of $C \rightarrow$ P vs. NP
can 3-SAT be computed in polynomial time?

$\downarrow$
poly(n) size, constant depth circuits can't compute parity

Switching Lemma f-DNF width w , δ -random restriction (baby)

$\Pr[f|_{\mathcal{R}} \text{ is not const}] \leq 5\delta w$, $3\delta w$ today's problems

$\Pr[DT(f|_{\mathcal{R}}) \geq k] \leq (5\delta w)^k$
 $\hookrightarrow$ depth & decision tree computing $f|_{\mathcal{R}}$.

$\hookrightarrow$ PSPACE vs. PH (polynomial-time hierarchy)

Lemma. $\epsilon = \Pr[DT(f|_{\mathcal{R}}) \geq k]$, then $(*)$
 f is 3ϵ -concentrated up to degree $3k/\delta$.

Note $\epsilon \geq \Pr[\deg(f|_{\mathcal{R}}) \geq k]$.

Corollary. If f is width- w DNF, ϵ -concentrated up to deg $O(w \log \frac{1}{\epsilon}) \rightarrow \|\hat{f}\|_2$

Lemma. f is width- w DNF,

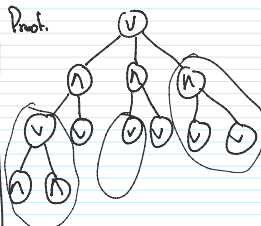
$$\sum_{|S| \leq k} |\hat{f}(S)| \leq 2 \cdot (20w)^k \rightarrow \|\hat{f}\|_1$$

f is ϵ -concentrated on $\mathcal{F}$, $|\mathcal{F}| \leq w^{O(w \log \frac{1}{\epsilon})}$
(goal is 2^w)

Thm. f size s , depth d width w circuit
 $\delta = \frac{1}{10w} \left(\frac{1}{10\epsilon}\right)^{d-2}$, $L = \log(\frac{2s}{\epsilon})$

$$\Pr[DT(f|_{\mathcal{R}}) \geq \log \frac{2s}{\epsilon}] \leq \epsilon$$

Proof.



Random restrictions in stages

Random restriction $\rightarrow$ decision tree $\rightarrow$ CNF
on DNF $\downarrow$
compress

Random restriction $\rightarrow$ decision tree $\rightarrow$ DNF
on CNF $\downarrow$
compress

$\Pr[\text{random restriction width} \geq w] \leq (\frac{3}{\epsilon})^w$
of AND

Problem 1 - Influence of DNF $\leq 2w$
CNF $\leq w$

Problem 2 - $3\delta w$

Problem 1 - Rows 1, 3

Problem 2 - Rows 2, 4