

Boolean Functions

1. Summary

1. DNF's

2. Random Restrictions

2. Q+A

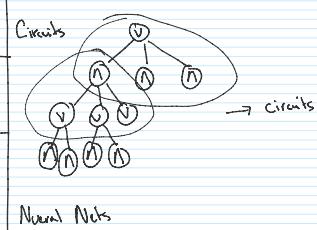
3. Problems

Announcements

- Solutions accessible

DNF - Formula OR of ANDs
width - max fanin of AND
size - fanin of OR

CNF - ANDs of ORs - 3-SAT



Prop. If f is a width- w DNF,
 $I[f] \leq 2w$, ϵ -concentrated up to degree $2w/\epsilon$

Can learn in time $\text{poly}(w)$

Theorem. If f is a size- s DNF, then $\forall g$,
 $\exists g$ computable by a width- $\log(s/\epsilon)$

DNF st. dist $(f, g) \leq \epsilon$

Corollary - size s , ϵ -concentrated up to
deg $O(\log(s/\epsilon)/\epsilon)$.

Mansour's conjecture - width- w ϵ -concentrated
on $2^{O(w)}$ coefficients

Mansour $w^{O(w)}$, $sO(\log s)$ Thus $I[f] \leq O(\log s)$ if size- s DNF δ - Random restriction

↳ for each input

+1 w/ prob $\frac{1-\delta}{2}$ -1 w/ prob $\frac{1+\delta}{2}$ variable w/ prob δ

$$\mathbb{E}[\hat{f}_{J_{12}}(S)] = \delta^{|S|} \cdot \hat{f}(S)$$

$$\mathbb{E}[\hat{f}_{J_{12}}(S)^2] = \sum_{U \subseteq [n]} \Pr[U \cap J = S] \cdot \hat{f}(U)^2 \\ = \sum_{S \subseteq U} \delta^{|S|} (1-\delta)^{|U-S|} \cdot \hat{f}(U)^2$$

$$\mathbb{E}[I[\hat{f}_{J_{12}}]] = \delta I[f]$$

DNF - OR of AND's

↓

 $\frac{1}{2}$ random restriction

for each and

Lemma A is AND, $\frac{1}{2}$ random restriction

$$\Pr[A_{J_{12}} \text{ is width } \geq w] \leq \left(\frac{3}{4}\right)^w$$

 $I[f] \leq 2w$ # of pairs, x, y s.t. $f(x) \neq f(y)$ $x+y$ differ in 1 coordinate

$$f(x) = \text{True}, f(y) = \text{False}$$

True → False only 1 of AND's can true
only w ways

$$e^{-\delta^2 \cdot C}$$

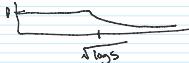
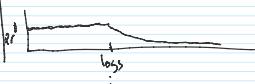
$$s \cdot e^{-\delta^2 \cdot C}$$

Union bound

$$\mathbb{E}[\max \text{ width of ANDs}] = \sum_{i=1}^w \Pr[\max \text{ width} \geq i]$$

Generic chaining

$$= \sum_{i=1}^w \min(1, s \left(\frac{3}{4}\right)^i)$$



Problem 1 - constructing a DNF for any function → before 2, 4 1, 5

Problem 2 - monotone functions 2, 4

Problem 3 - $\text{Inf}[f_{J_{12}}] = \delta \text{Inf}[f]$ w/o Fourier Analysis 3, 7 > RoomsProblem 4 - $X_{[n]}$, XOR/Parity 4, 8AC₀ - constant depth polynomial-size circuits, Parity & AC₀