

Be The Leader

Alg 0: Be the Leader (BTL)

- let $V_j^i = \sum_{r=1}^i v_j^r$
- in round i choose:
 $j^i = \arg\max_j V_j^i$

(Recall: FTL $j^i = \arg\max_j V_j^{i-1}$)

Example

$k=2$ actions

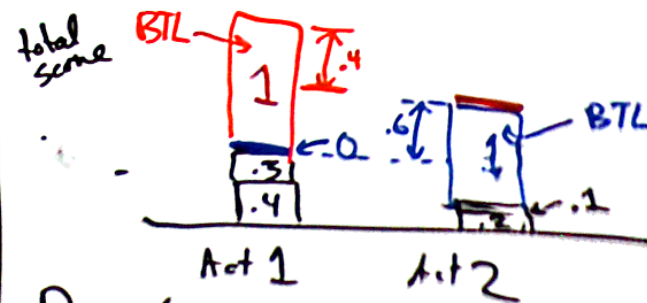
	1	2	3	4
Action 1	.4	.3	0	1
Action 2	.2	.1	1	0
BTL total	.4	.7	1.7	2.7
OPT total	.4	.7	1.3	1.7

best action in hindsight

Thm: $BTL \geq OPT$.

"change in BTL in each round"

$\geq$
 "change in OPT in each round"



Proof:

- let $OPT_i =$ best in hindsight after i rounds
- let $opt_i = OPT_i - OPT_{i-1}$
- let $btl_i =$ passoff BTL in round i
- claim: $btl_i \geq opt_i$
- $btl_i =$ passoff of best in hindsight action $\geq$ increase in best in hindsight

btl_i = payoff of best in hindsight action
 $\geq$ increase in total payoff of best in hindsight.

$= opt_i$

$BTL = \sum_i btl_i \geq \sum_i opt_i = OPT$

Follow The Perturbed Leader

Alg 2 : Follow the perturbed leader (FTPL)

• learning rate ϵ

• hallucinate:

$V_j^0 = h \times$ "number of tails of ϵ -biased coin that flip in a row"

• let $V_j^i = \sum_{r=1}^i V_j^r$

• in round i choose action

$j^i = \arg\max_j (V_j^0 + V_j^{i-1})$

Example:

	0	1	2	3	4	5
Action 1	2	$\boxed{1/2}$	$\boxed{0}$	1	$\boxed{0}$	$\boxed{1}$
Action 2	3	$\boxed{0}$	$\boxed{1}$	$\boxed{0}$	$\boxed{1}$	$\boxed{0}$

$OPT : 2.5 \rightarrow n/2$

$FTPL : 2 \rightarrow n/2$

$FTL : 1/2 \rightarrow 1/2$

Thm : payoffs in $[0, h]$

$E[FTPL] \geq (1-\epsilon)OPT - O(h \sqrt{\log k})$

Cor : in n rounds with payoffs in $[0, h]$

$E(\text{Regret}_n(\text{FTPL})) \leq 2h \sqrt{\frac{\log k}{n}}$

Proof of Cor : same as for EW.

Q: why does FTPL work?

A:

1) stability: $FTPL \approx BTPL$

2) small perturbation: $BTPL \approx OPT$

Def: Be the Perturbed Leader (BTPL)

$$j^i = \arg \max_j (v_j^0 + V_j^i)$$

Lemma 1 (stability): $FTPL \geq (1-\epsilon) BTPL$

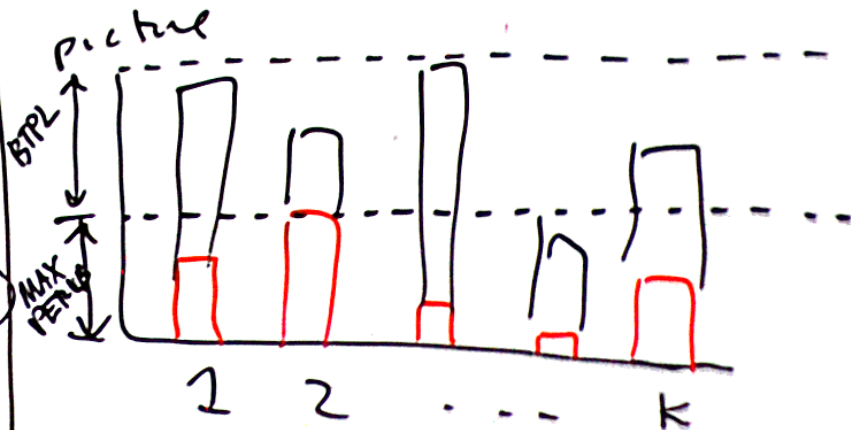
Lemma 2 (Small Perturb):

$$BTPL \geq OPT - O\left(\frac{1}{\epsilon} \log k\right)$$

Proof of Thm:

combine lemmas.

Proof of Lemma 2:



$$Thm \Rightarrow \widetilde{BTPL} \geq \widetilde{OPT} \geq OPT$$

$\widetilde{BTPL} = BTPL + MAXPERTURB$

$$\Rightarrow BTPL \geq OPT - MAXPERTURB$$

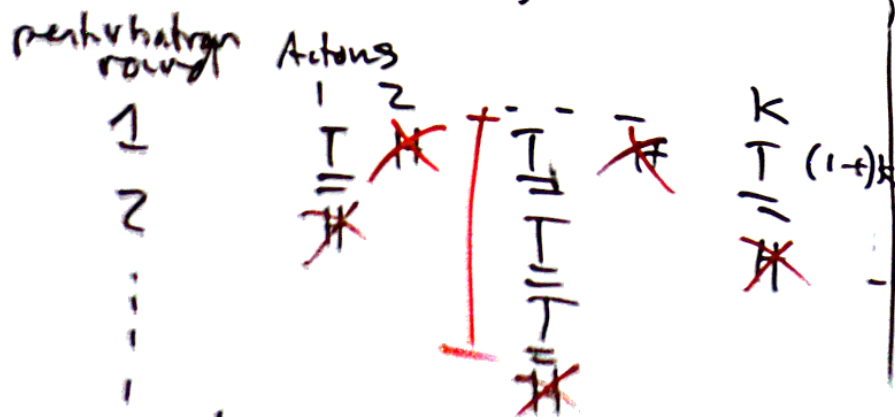
$$claim: [MAXPERTURB] = O\left(\frac{1}{\epsilon} \log k\right)$$

MAXPERTURB = #rounds until no actions left.

$$e.g. \epsilon = 1/2 \Rightarrow \log k$$

- $V_j^0 = h \times$ "number of ϵ bias coins that flip tail in a row"
- $= h \times$ Geometric Random Variable with $p = 1 - \epsilon$.

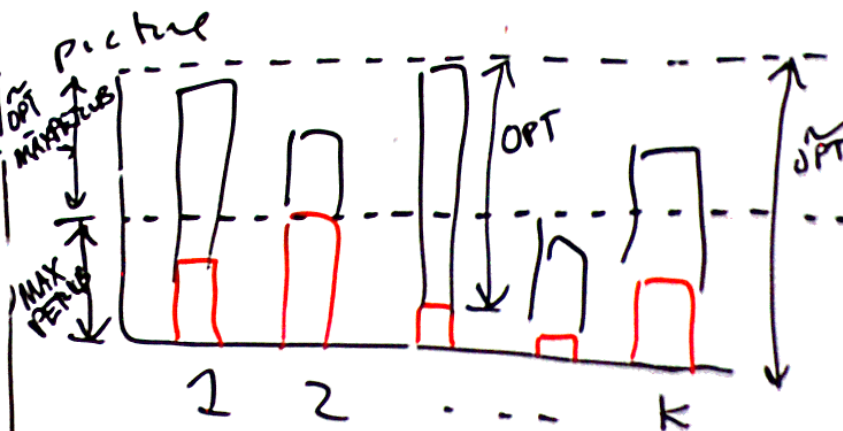
$$\text{MAXPERTURB} = \max_j V_j^0$$



"log $\frac{1}{1-\epsilon}$ k"

$\frac{1}{\ln \frac{1}{1-\epsilon}} \ln k \rightarrow \ln \frac{1}{1-\epsilon}$
 $= -\ln(1-\epsilon)$
 $\approx \epsilon$

Proof of Lemma 2:



$$\text{Thm} \Rightarrow \widetilde{\text{BTPL}} \geq \widetilde{\text{OPT}} \geq \text{OPT}$$

$$\text{BTPL} + \text{MAXPERTURB}$$

$$\Rightarrow \text{BTPL} \geq \text{OPT} - \text{MAXPERTURB}$$

claim: $\text{MAXPERTURB} = O\left(\frac{1}{\epsilon} \log k\right)$

MAXPERTURB = #rounds until no actions left

e.g. $\epsilon = \frac{1}{2} \Rightarrow \log k$