

Comp Sci 212

1. Markov's Inequality

2. Chebyshev's Inequality

Announcements

Midterm grades out

Getting to know you meetings

Tails bounds

R random variable

$$\Pr[R \geq x] \leq ?$$

- If R represent the running time of a random algorithm (or the error), want $\Pr[R \geq x]$ to be small.

Markov's Inequality

What tails bounds bounds can you obtain from just the expected value?

Ex. In a room of 100 people avg age 40

S = set of people

$$\Pr[w] = \frac{1}{100} \text{ for each } w \in S$$

 $R: S \rightarrow \mathbb{R}$, $R(w)$ age of person w

$$E[R] = 40 = \sum_{w \in S} R(w) \cdot \Pr[w]$$

What is the largest $\Pr[R \geq 100]$ can be?
proportion of people ≥ 100 years

Not everyone can be at least 100

Even $\Pr[R \geq 100] = \frac{1}{2}$ is impossible
↓
worst case 50 people 100

50 people 0, \rightarrow avg. 50

$$\Pr[R \geq 100] = \frac{1}{4}$$
 is possible

25 people be 100 \rightarrow avg. 40

75 people be 20

$$\text{largest is } \Pr[R \geq 100] = \frac{4}{10} = \frac{E[R]}{100}$$

40 people 100 \rightarrow avg. 40

60 people 0

make everyone either 100 or 0.

Markov's Inequality

If R is a non-negative random variable

 $x > 0$, then

$$\Pr[R \geq x] \leq \frac{E[R]}{x}$$

Proof. Eough to prove

$$E[R] \geq x \Pr[R \geq x]$$

$$\begin{aligned} \sum_{w \in S} R(w) \Pr[w] &\stackrel{\downarrow}{\rightarrow} \sum_{w \in S} \Pr[w] \cdot x \\ \text{at least } 0 & \\ \sum_{w \in S} R(w) \Pr[w] &= \sum_{w \in S} R(w) \Pr[w] + \sum_{w \in S, R(w) \geq x} R(w) \Pr[w] \\ &\geq \sum_{w \in S, R(w) \geq x} R(w) \Pr[w] \quad (R(w) \geq x) \\ &\geq \sum_{w \in S, R(w) \geq x} x \Pr[w] \quad (\Pr[w] \geq 0) \end{aligned}$$

Coins - $S = \{H, T\}^{100}$ $R: S \rightarrow \mathbb{R}$, $R(w) = \# \text{ of heads in } w$ Uniform dist., sticky, $E[R] = 50$

$$\Pr[R \geq 75] \leq \frac{E[R]}{75} \leq \frac{2}{3}$$

(by Markov's inequality)

Shaky coin: $\Pr[R \geq 75] = \frac{1}{2}$ ↓ all heads up $\frac{1}{2}$ all tails up $\frac{1}{2}$ Uniform dist.: $\Pr[R \geq 75] = \text{really small}$

(example Markov's inequality doesn't give a good bound)

Birthday Paradox: $D = \text{set of days, } |D|=365$ $S = D^n$ (sample space), uniform dist. $R: D^n \rightarrow \mathbb{R}$, $R(w) = \# \text{ of pairs of people w/ same birthday}$ $E_{ij} = \text{event that person } i \neq \text{person } j \text{ have same birthday}$

$$\Pr[E_{ij}] = \frac{1}{365}$$

$$R = \sum_{\text{pairs } ij} \Pr_{ij} \rightarrow E[R] = \sum_{\text{pairs } ij} \Pr[E_{ij}]$$

$$= \frac{\# \text{ of pairs}}{365} = \frac{n(n-1)}{730}$$

If $n=27$, $E[R] \approx 1$

$$\Pr[R \geq 10] \leq \frac{E[R]}{10} \approx \frac{1}{10}$$

(by Markov's inequality)

Back to room (avg age 40)

Assume that everyone is at least 20 years old

40 people above 100 is now impossible

40 people 100

60 people 20 avg 52

largest possible is $\Pr[R \geq 100] \leq \frac{25}{100}$

25 people 100 years

75 people 20 years

$$\begin{aligned} E[R] &= 40 \\ b &= 20 \\ \frac{40-20}{100-20} &= \frac{20}{80} = \frac{1}{4} \end{aligned}$$

Corollary: If $R(w) \geq b$ for all w , $x > b$

$$\Pr[R \geq x] \leq \frac{E[R]-b}{x-b}$$

(b=0, just Markov's inequality)

Proof. Let $Y(w) = R(w) - b$ $Y(w)$ is always non-negative

$$E[Y] \leq E[R] - b$$

$$\Pr[R \geq x] = \Pr[Y \geq x-b] \leq \frac{E[Y]}{x-b} = \frac{E[R]-b}{x-b}$$

(Markov's inequality)

