

Comp Sci 212

What can we say about a random variable, based on just its expected value?

$$P_i[R \geq x] \leq ?$$

random variable number

- If R represents the running time of a random algorithm, want $\Pr[R \geq x]$ to be small.

Def. Expected Value Sample space S , Prob. list P

$$E[R] = \sum_{w \in S} R(s) \cdot P_r[w] \quad R: S \rightarrow \mathbb{R}$$

Ex. In a room with 100 people, average age is 40. How many people can be at least 100 years old?

S = set of people

$\Pr[w] = \frac{1}{100}$ for all $w \in S$, uniform

$R: S \rightarrow \mathbb{R} \quad R(w) = \text{age of person } w$

$$E[R] = \sum_{w \in S} \frac{R(w)}{100} = 40$$

What is the largest $\Pr[R \geq 100]$ can be?

- Not everyone can be at least 100 $\rightarrow P_r[R \geq 100] < 1$
 - average age would be at least 100
- Not even 50 people can be over 100 $\rightarrow P_r[R \geq 100] < \frac{1}{2}$
 - if 50 people are 100, 50 are 0, average is 50
- 25 people is possible to be 100, avg. is 25
 75 people 20 years, avg. is 40

- largest 40 people

40 people 100 years old
60 people 0 years old $\rightarrow$ avg is 40

$$P_r[R \geq 100] \leq \frac{40}{100} = \frac{E[R]}{100}$$

Thm. Markov's Inequality

If R is a non-negative random variable ($R(\omega) \geq 0$ for all $\omega \in S$), $x > 0$

$$P[R \geq x] \leq \frac{E[R]}{x}$$

Proof. Enough to prove $x \cdot \Pr[R \geq x] \leq E[R]$

$$\sum_{\substack{w \in S \\ \text{s.t. } R(w) \geq x}} R(w) \cdot P_R(w)$$

$$\begin{aligned} \sum_{w \in S} R(w) P_r[w] &= \sum_{\substack{w \in S \text{ s.t.} \\ R(w) \geq x}} R(w) P_r[w] + \underbrace{\sum_{\substack{w \in S \text{ s.t.} \\ R(w) < x}} R(w) P_r[w]}_{\text{at least } 0} \\ &\geq \sum_{\substack{w \in S \text{ s.t.} \\ R(w) \geq x}} R(w) P_r[w] \\ &\geq \sum_{\substack{w \in S \text{ s.t.} \\ R(w) \geq x}} x \cdot P_r[w] = x \cdot P_r[R \geq x] \end{aligned}$$

- Assume that everyone is at least 20 years old in our room.

- Now 40 people over 100 is impossible

- 40 people 100 years avg. 52

- 60 people 20 years

- largest # is 25 people $\rightarrow P_c[R \geq 40] \leq \frac{1}{4}$
 25 people 100 years old
 75 people 20 years old avg. 40

$$\rightarrow \Pr[E_{i,j}] = \frac{1}{|D|}$$

$E_{i,j}$ = event that person i & j have the same birthday, $R = \sum_{\substack{i,j \text{ st.} \\ i < j}} 1_{E_{i,j}}$, $E[R] = \sum_{\substack{i,j \text{ st.} \\ i < j}} P_r[E_{i,j}]$

$$= \sum_{\substack{i, j \in \mathcal{V} \\ i < j}} \frac{1}{|\mathcal{N}|} = \frac{n \cdot (n-1)}{2|\mathcal{N}|} \approx 1$$

Birthdays Paradox - $D =$ set of possible days
 n people in a room, $S = D \times D \times \dots \times D$, uniform dist.

$R: S \rightarrow R$, $R = \#$ of pairs of people with the same birthday

If $n \geq 2\sqrt{101} \rightarrow \Pr[R \geq 1] \geq 1 - \frac{1}{2}$, but what about $\Pr[R \geq 10]$?

$\hookrightarrow \leq \frac{1}{5}$ by Markov's inequality

Corollary - If $R(\omega) \geq b$ for all ω ,

$$Pr[R \geq x] \leq \frac{E[R] - b}{x - b}$$

Proof. Let $Y(\omega) = R(\omega) - b$

γ is always non-negative

$$E[Y] = E[R - b] = E[R] - b$$

$$\Pr[R \geq x] = \Pr[Y \geq x-b] = \frac{E[Y]}{x-b} = \frac{E[R]-b}{x-b}$$

Room examples show that Markov's inequality and Corollary are tight.

$S = \{H, T\}$, $S = \underbrace{S' \times \dots \times S'}_{100 \text{ times}}$ uniform
sticky - all heads
all tails

$R: S \rightarrow \mathbb{R}$ $R(s) = \# \text{ of heads}$. $E[R] = 50$
both cases

By Markov's inequality $P[R \geq 75] \leq \frac{50}{75} \leq \frac{2}{3}$

Right answer - sticky $\Pr[R \geq 75] = \frac{1}{2}$
uniform $\Pr[R \geq 75] = \text{really small}$